## Field dynamics: a new tool to boost mixing results

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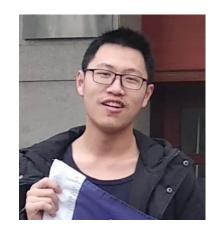
Joint work with:



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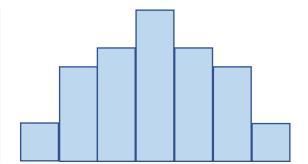
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Summer school at UCSB, Santa Barbara, CA, US, 12th August 2022

# Sampling, counting and phase transition

Boolean variables set V, weight function  $w: \{-, +\}^V \to \mathbb{R}_{\geq 0}$  joint distribution  $\mu$ :

$$\forall X = (X_v)_{v \in V} \in \{-, +\}^V, \qquad \mu(X) \propto w(X)$$

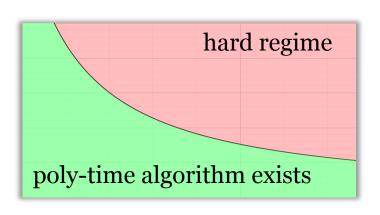


## Sampling problem

Draw (approximate) random samples from distribution  $\mu$ 

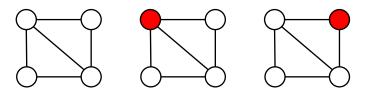
#### Goal:

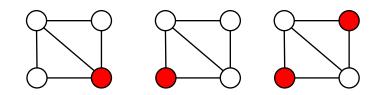
Prove *optimal* mixing results up to the computational phase transition threshold



## Example: Hardcore model

- graph G = (V, E), parameters  $\lambda$ ;
- Gibbs distribution  $\mu$ :  $\forall$  independent set  $I \subseteq V$ ,  $\mu(I) \propto \lambda^{|I|}$ .
- Equivalent state space of  $\mu$ :  $\{-,+\}^V = \{\text{occupied}, \text{unoccupied}\}^V$





#### **Computational phase transition**

- $\lambda < \lambda_c(\Delta)$ : poly-time algorithm for sampling [Weitzo6]
- $\lambda > \lambda_c(\Delta)$ : no poly-time algorithm unless NP = RP [Sly10]

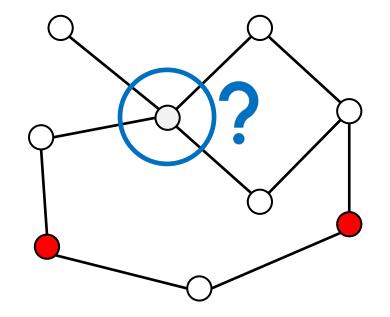
$$\lambda_c(\Delta) = \frac{(\Delta - 1)^{(\Delta - 1)}}{(\Delta - 2)^{\Delta}}$$
$$\approx \frac{e}{\Delta}$$

# Glauber dynamics for hardcore model

Start from an arbitrary independent set *X*;

For each transition step do

- Lazy w.p.  $\frac{1}{2}$ , otherwise do as follows:
- Pick a vertex *v* uniformly at random;
- If  $X_u = -$  for all neighbors u then  $X_v = \begin{cases} + & \text{w. p. } \lambda/(1+\lambda) \\ & \text{w. p. } 1/(1+\lambda) \end{cases}$
- Else  $X_v \leftarrow -$



Mixing time:  $T_{\text{mix}} = \max_{X_0 \in \Omega} \min \left\{ t \mid d_{TV}(X_t, \mu) \leq \frac{1}{4e} \right\}$ ,

 $d_{TV}(X_t, \mu)$ : the *total variation distance* between  $X_t$  and  $\mu$ .

## Previous works

Work	Condition	<b>Mixing Time</b>
Dobrushin 1970	$\lambda \le \frac{1-\delta}{\Delta-1}$	$O\left(\frac{1}{\delta}n\log n\right)$
Luby, Vigoda, 1999	$\lambda \le \frac{2(1-\delta)}{\Delta - 2}$	$O\left(\frac{1}{\delta}n\log n\right)$
Efthymiou et al, 2016	$\lambda \leq (1 - \delta)\lambda_c(\Delta)$ $\Delta \geq \Delta_0(\delta)$ , girth $\geq 7$	$O\left(\frac{1}{\delta}n\log n\right)$

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Anari, Liu, Oveis Gharan, 2020 improved by Chen, Liu, Vigoda, 2020	$\lambda \le (1 - \delta)\lambda_c(\Delta)$	$n^{O(1/\delta)}$
Chen, Liu, Vigoda, 2021	$\lambda \le (1 - \delta)\lambda_c(\Delta)$	$\Delta^{O(\Delta^2/\delta)} n \log n$

## **Open question:**

Can we prove the fast (optimal) mixing for all degrees?

## Mixing time of Glauber dynamics when $\lambda \leq (1 - \delta)\lambda_{\mathcal{C}}$

Work	Mixing Time	Technique
Anari, Liu, Oveis Gharan, 2020 improved by Chen, Liu, Vigoda, 2020	$n^{O(1/\delta)}$	Spectral Independence (SI)
Chen, Liu, Vigoda, 2021	$\Delta^{O(\Delta^2/\delta)} n \log n$	
Chen, F., Yin, Zhang, 2021	$e^{O(1/\delta)}n^2\log n$	SI & Field Dynamics

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3 Chen, Eldan, 2022	$e^{O(1/\delta)}n\log n$	Localization Scheme

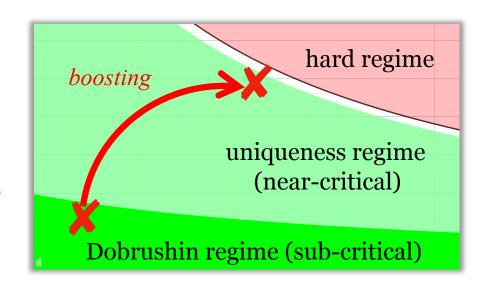
Zongchen Chen

<sup>2</sup> Xiaoyu Chen

#### Hardcore model in uniqueness regime

If  $\lambda$  is **close** to  $\lambda_c(\Delta)$ , e.g.,  $\lambda = 0.999\lambda_c$  (**near-critical**) analyzing mixing time is **hard** 

• If  $\lambda$  is *far-away* from  $\lambda_c(\Delta)$ , e.g.,  $\lambda \leq 0.1\lambda_c$  (sub-critical) analyzing mixing time is *easy* 



## **Boosting Theorem**

Boosting mixing results from sub-critical regime to near-critical regime

- Boost spectral gap of Glauber dynamics [CFYZ21]
- Boost modified log-Sobolev constant of Glauber dynamics [CFYZ22]

Proved by a new Markov chain: field dynamics

## Revisit Chen-Liu-Vigoda's technique

Simpler task: poly-time sampling algorithm

**Input**: hardcore model with  $\lambda \leq (1 - \delta)\lambda_c(\Delta)$  and  $\Delta$  can be unbounded;

**Output:** random sample *X* s.t.  $d_{TV}(X, \mu) = \frac{1}{\text{poly}(n)}$ .

## $\theta$ -fractional block dynamics

Parameter:  $\theta \in (0,1)$ 

Initialization: arbitrary  $X \in \{-, +\}^V$ 

Update: for each t = 1 to T

- pick  $S \subseteq V$  with  $|S| = \theta n$  u.a.r.;
- $X_S \sim \mu_S(\cdot | X_{V \setminus S});$

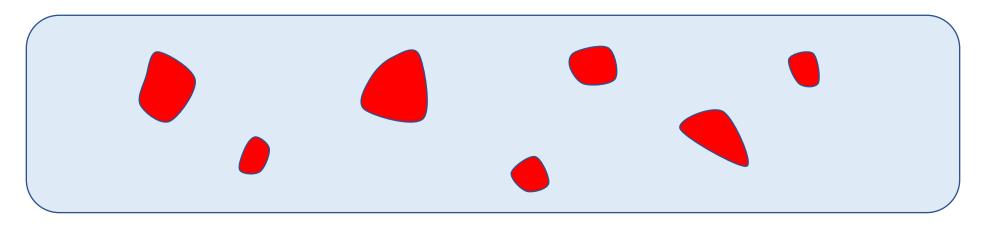
## **Mixing result [CLV21]**

$$\lambda_{\rm gap} \ge \theta^{O(1/\delta)}$$

$$d_{TV}(\mu, X) \le \frac{1}{\text{poly}(n)} \text{ if } T = \left(\frac{1}{\theta}\right)^{O(1/\delta)} n \log n$$

**Question**: how to *efficiently* simulate the transition step  $X_S \sim \mu_S(\cdot | X_{V \setminus S})$ ?

Update step  $X_S \sim \mu_S(\cdot | X_{V \setminus S})$ : sample from hardcore model  $(G[S], \lambda)$  with boundary condition  $X_{V \setminus S}$ 



Observation [Chen, Liu and Vigoda, 2021]

If  $\theta = O\left(\frac{1}{\Delta}\right)$ , then w.h.p., G[S] is a set of small connected components

$$\theta = 0(1/\Delta)$$
 fractional block dynamics 
$$T = \Delta^{0(1/\delta)} n \log n \text{ steps}$$

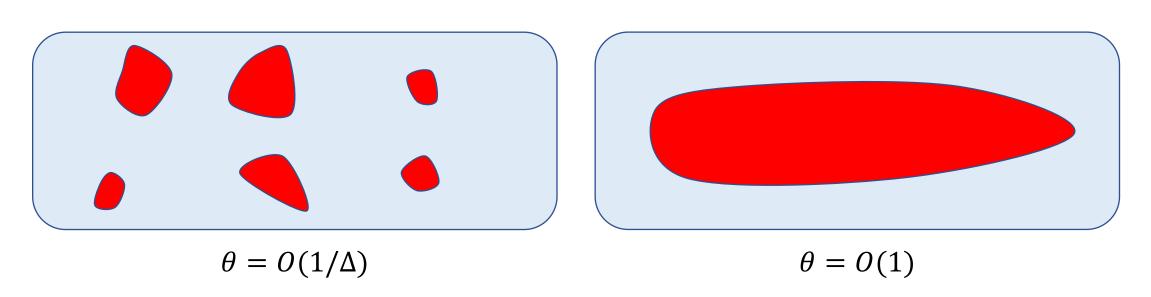
Simulation of each step brute force on each component Expected cost = O(n)

Total expected running time of the algorithm:  $\Delta^{O(1/\delta)} n^2 \log n$ 

Natural idea: set 
$$\theta = \frac{1}{100}$$



Mixing time  $T = 2^{O(1/\delta)} n \log n = O_{\delta}(n \log n)$ 



**Issue**: how to sample from hardcore model (G[S],  $\lambda$ ) with boundary condition  $X_{V \setminus S}$ ?

### Observation [Chen, F. Yin and Zhang, 2021]

The maximum degree of G[S] can be small.

For any 
$$v \in S$$
,  

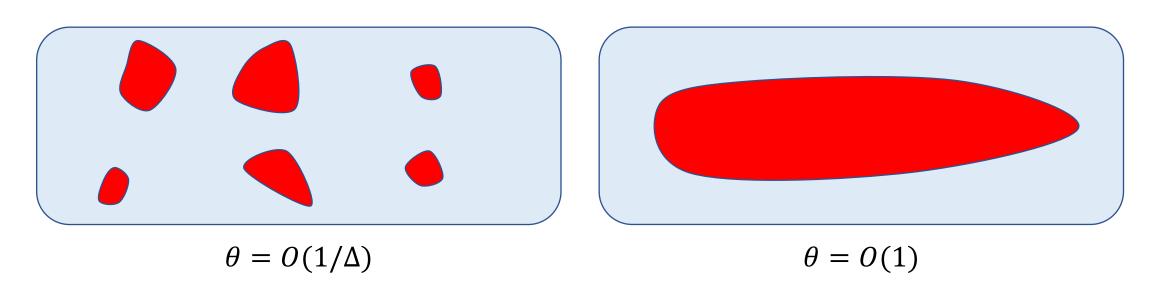
$$\mathbb{E}[\text{degree of } v \text{ in } G[S]] \approx \theta \deg_G(v)$$

$$= \frac{\deg_G(v)}{100}$$

Natural idea: set 
$$\theta = \frac{1}{100}$$



Mixing time  $T = 2^{O(1/\delta)} n \log n = O_{\delta}(n \log n)$ 



**Issue**: how to sample from hardcore model  $(G[S], \lambda)$  with boundary condition  $X_{V \setminus S}$ ?

$$\lambda \leq \lambda_c(\Delta_G) \approx \frac{e}{\Delta_G}$$

uniqueness condition in  $(\lambda, \Delta_G)$ 

If we can show

$$\Delta(G[S]) \ll \Delta_G$$

the  $(G[S], \lambda)$  is **easy** to sample from

**Case 1**: the maximum degree  $\Delta_G$  of original graph G satisfies  $\Delta_G \geq 100 \log n$ 

• By **concentration**, for any  $v \in S$ , expected degree  $\leq \Delta_G/100$ ,

$$\Pr\left[\text{degree of } v \text{ in } G[S] \leq \frac{\Delta_G}{10}\right] \geq 1 - \frac{1}{n^{10}}$$

• Bound a **union bound**, w.h.p.  $(\text{prob} \ge 1 - \frac{1}{n^7})$ 

In every transition step, the maximum degree of G[S] is at most  $\frac{\Delta_G}{10}$ 

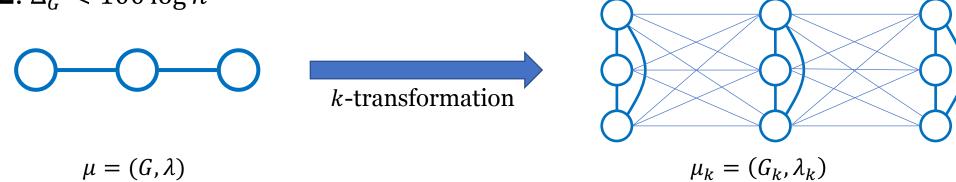
• The hardcore model (G[S],  $\lambda$ ) satisfies **Dobrushin's condition** simulate the Glauber dynamics for  $O(n \log n)$  steps.

$$\theta = 1/100$$
-fractional block dynamics  $T = 2^{O(1/\delta)} n \log n$  steps

simulation cost of each step  $O(n \log n)$ 

Total running time of the algorithm:  $2^{O(1/\delta)}n^2 \log^2 n$ 

**Case 2**:  $\Delta_G < 100 \log n$ 

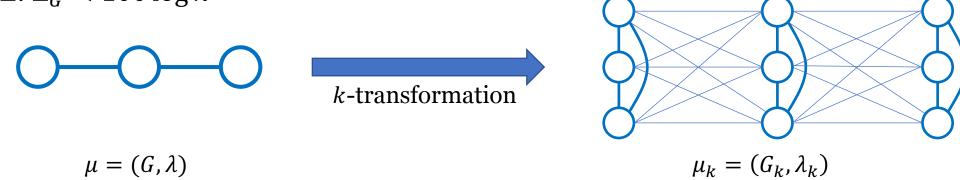


- Graph  $G_k: v \in V \longrightarrow \text{ size } k \text{ clique } C_v; \quad \{u, v\} \in E \longrightarrow \text{ connect } C_u \text{ and } C_v;$
- Parameter:  $\lambda_k = \lambda/k$ ;

#### **Properties of the** *k***-transformation:**

- $\mu_k$  is  $O(1/\delta)$  spectrally independence  $\longrightarrow$  block dynamics on  $\mu_k$  is rapid mixing [CLV21]
- if  $k = \Omega(\log n)$ , the **max degree** of  $G_k$  is **large**  $\longrightarrow \mu_k = (G_k, \lambda_k)$  is in **Case 1**
- if  $X \sim \mu_k$ , then  $X' = f_k(X) \sim \mu$

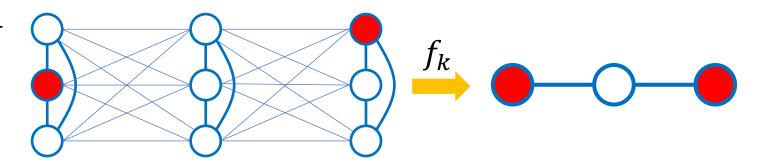
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- if  $X \sim \mu_k$ , then  $X' = f_k(X) \sim \mu$ 
  - $X'_v = + \text{ if } \exists u \in C_v \text{ s.t. } X_u = +$
  - $X'_v = -if \ \forall u \in C_v, \ X_u = -$



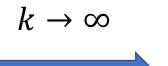
#### Algorithm for sampling from the hardcore model

**Input**: graph G = (V, E) and parameter  $\lambda \leq (1 - \delta)\lambda_c(\Delta)$ 

- $\triangleright$  apply  $k = \Omega(\log n)$ -transformation to get  $\mu_k$ ;
- > simulate  $\left(\theta = \frac{1}{100}\right)$ -fractional block dynamics  $(X_t)_{t=0}^T$  on  $\mu_k$ :
  - $T = 2^{O(1/\delta)}(nk)^2 \log(nk)$ ;
  - every transition is simulated by an  $O(nk \log(nk))$ -step **Glauber dynamics**
- $\triangleright$  output  $f_k(X_T)$

$$\left(f_k(X_t)\right)_{t=0}^T$$

Apply the mapping  $f_k$  on every step of block dynamics



Field Dynamics

New Markov chain for sampling from  $\mu = \text{Hardcore}(G, \lambda)$ 

## Field Dynamics

**Input**: hardcore model  $\mu = (G, \lambda)$ , a parameter  $\theta \in (0,1)$ 

Start from an arbitrary feasible configuration  $X \in \{-, +\}^V$ 

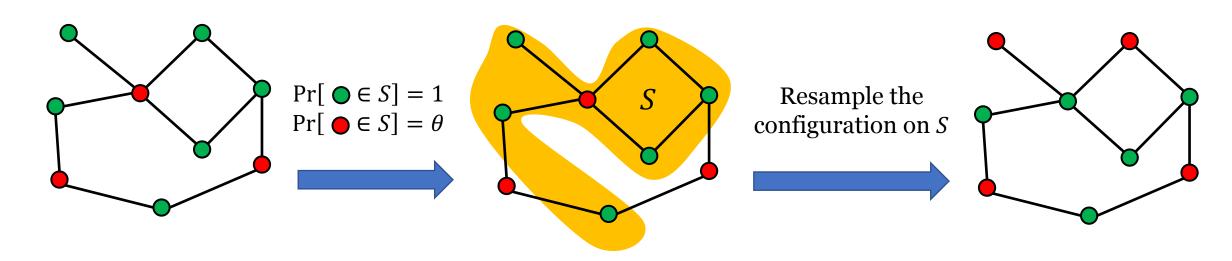
For each t from 1 to T do

**Down-Walk** • Construct  $S \subseteq V$  be selecting each  $v \in V$  independently with probability

$$p_v = \begin{cases} 1 & \text{if } X_v = -\\ \theta & \text{if } X_v = + \end{cases}$$

**Up-Walk** • Resample  $X_S \sim \pi_S(\cdot | X_{V \setminus S})$ 

 $\pi$ : hardcore model  $(G, \theta \lambda)$ 



### Field Dynamics

**Input**: a **general distribution**  $\mu$  over  $\{-1, +1\}^V$ , a parameter  $\theta \in (0, 1)$ 

Start from an arbitrary feasible configuration  $X \in \{-, +\}^V$ 

**For** each *t* from 1 to *T* **do** 

• Construct  $S \subseteq V$  be selecting each  $v \in V$  independently with probability

$$p_v = \begin{cases} 1 & \text{if } X_v = -\\ \theta & \text{if } X_v = + \end{cases}$$

• Resample  $X_S \sim \pi_S(\cdot | X_{V \setminus S})$ 

$$\forall \sigma \in \{-,+\}^V, \qquad \pi(\sigma) \propto \mu(\sigma) \prod_{v \in V: \sigma_v = +} \theta$$

 $\pi$ : distribution  $\mu$  with external field  $\theta$ 

## Field Dynamics

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• Resample  $X_S \sim \pi_S(\cdot | X_{V \setminus S})$ 

**Proposition** (Field Dynamics): for any  $\theta \in (0,1)$ 

Field dynamics has the unique stationary distribution  $\mu$ .

(irreducible, aperiodic and reversible)

# Spectral gap of Glauber dynamics

### Mixing lemma

If 
$$\lambda \leq (1 - \delta)\lambda_c(\Delta)$$
, for any  $\theta \in (0,1)$   

$$\operatorname{Gap}_{\text{field}}(\lambda, \theta) \geq \theta^{O(1/\delta)}$$

## Comparison lemma

For any  $\lambda \geq 0$ , for any  $\theta \in (0,1)$ 

$$Gap_{Glauber}(\lambda) \ge Gap_{field}(\lambda, \theta) \cdot Gap_{Glauber}(\theta\lambda)$$

## **Boosting theorem**

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Proved by a calculation

$$\theta = \frac{1}{10}$$

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$$Gap_{Glauber}(\theta\lambda) = \Omega\left(\frac{1}{n}\right)$$

$$\theta\lambda \in Dobrushin's regime$$

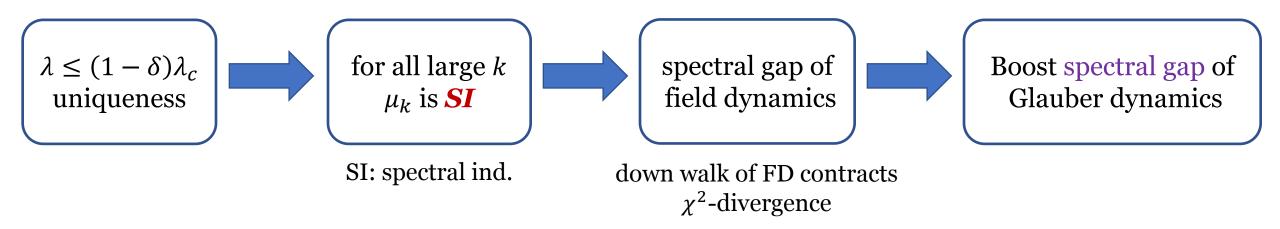
$$Gap_{Glauber}(\lambda) = \Omega_{\delta}\left(\frac{1}{n}\right)$$

$$T_{mix} = O_{\delta}(n^2 \log n)$$

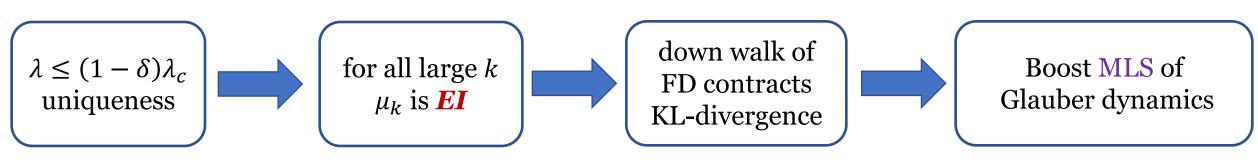


$$Gap_{Glauber}(\lambda) = \Omega_{\delta} \left(\frac{1}{n}\right)$$
$$T_{mix} = O_{\delta}(n^2 \log n)$$

## **Boosting spectral gap of Glauber dynamics**



## Boosting modified log-Sobolev constant of Glauber dynamics



EI: entropic ind.

$$MLS_{Glauber}(\lambda) = \Omega_{\delta}\left(\frac{1}{n}\right)$$
 and  $T_{mix} = O_{\delta}(n \log n)$ 

#### **Summary**

- Hardcore model in the uniqueness regime
  - Optimal spectral gap and  $O(n^2 \log n)$  mixing time
  - Optimal modified log-Sobolev constant and  $O(n \log n)$  mixing time
- General distributions
  - Complete SI ------ boost spectral gap
  - Complete SI + marginal ratio bound —— boost MLS constant Thank you!
- A new Markov chain field dynamics

#### Open problem

- More applications of field dynamics
  - Algorithmic applications (e.g., random cluster model [Chen and Zhang 2022])
- Extend our technique to *general distributions* beyond the Boolean domain i.e., q-coloring